





Information Theoretic Model Validation by Approximate Optimization

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Overview

- Motivation of information theory for optimization
- Approximation capacity of a cost function
- Examples
 - Binary symmetric channel
 - Cluster validation
 - role mining for role-based access control (RBAC)
 - Robust SVD

Optimization approach to pattern recognition

- Given: data $\mathbf{X} \in \mathcal{X}$ in data (input) space \mathcal{X}
- Goal: Learn structure from data, i.e., interpret data relative to a hypothesis class
- Hypothesis class C with hypotheses (solutions)

$$egin{array}{rcl} c & \colon & \mathcal{X} & o & \mathbb{K} & ext{(e.g., } \mathbb{B}^n ext{ or } \{1, \dots, k\}^n) \ & \mathbf{X} & \mapsto & c(\mathbf{X}) \end{array}$$

• Cost function to define a partial order on \mathcal{C}

Pattern recognition and modeling

- Given are data and interpretations of these data,
 i.e., hypotheses.
- Modeling is (partial) ranking of the hypotheses encoded as data dependent costs.
 - Good/poor hypotheses have low/high costs
 - Optimal hypotheses minimize costs and are random variables.
- ⇒Search for hypotheses that have low costs on future data, i.e. **generalize** well!

Coding & pattern recognition with noisy data

IT: Space of strings is PR: Hypothesis class is partitioned by code vectors

partitioned by code problems



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Coding by Code Problems

- Idea: define a code by transforming a given optimization problem
- \Rightarrow codebook $\mathbb{T} = \{\tau_i \in \mathcal{T} : 1 \leq i \leq 2^{n\rho}\}$ with transformation set

$$\mathcal{T} = \{ \tau : R(c, \mathbf{X}) = R(\tau \circ c, \tau_{\mathbf{X}} \circ \mathbf{X}) \}$$

- Combinatorial optimization: permutation of vertices in a graph
- Identifiable transformations T are <u>messages</u>!

Asymptotical error-free communication $\lim_{n\to\infty} P(\hat{\tau} \neq \tau_s | \tau_s) = 0 \text{ is possible if } \dots$

• ... mutual information $\mathcal{I}_{\beta}(\tau_s, \hat{\tau})$ is bounded by

$$\rho < \mathcal{I}_{\beta}(\tau_{s}, \hat{\tau}) \equiv \frac{1}{n} \log_{2} \frac{|\mathcal{T}| \ Z_{\beta}^{(1\&2)}}{Z_{\beta}^{(1)} \ Z_{\beta}^{(2)}} \\ = \frac{1}{n} \left(\log_{2} \frac{|\mathcal{T}|}{Z_{\beta}^{(1)}} + \log_{2} \frac{|\mathcal{C}^{(2)}|}{Z_{\beta}^{(2)}} - \log_{2} \frac{|\mathcal{C}^{(2)}|}{Z_{\beta}^{(1\&2)}} \right) \\ \text{Bound calculation involves} \qquad Z_{\beta}^{(\nu)} = \sum_{c \in \mathcal{C}(\mathbf{X}^{(\nu)})} \exp(-\beta R(c, \mathbf{X}^{(\nu)}), \ \nu = 1, 2) \\ \text{and joint costs} \qquad Z_{\beta}^{(1\&2)} = \sum_{c \in \mathcal{C}(\mathbf{X}^{(1)})} \exp(-\beta (R(c, \mathbf{X}^{(1)}) + R(c, \mathbf{X}^{(2)})) \right) \\ \end{array}$$

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Model Selection by Maximization of Approximation Capacity



Maximize channel capacity w.r.t. approximation quality β, topology and metric of solution space, cost function *R*(.,.)

ASC for binary channel consistent with Shannon information theory

 $\xi^{(1)}$

}

- Hypothesis class: set of binary strings $\xi^{(1)}, \xi^{(2)} \in \{-1, 1\}^n$
- Communication:
- Costs of string s: Hamming distance

$$R(s,\xi^{(1)}) = \sum_{i=1}^{n} \mathbb{I}_{\{s_i \neq \xi_i^{(1)}\}}$$

• Mutual information: $\mathcal{I}_{\beta} = \ln 2 + (1 - \delta) \ln \cosh \beta - \ln(\cosh \beta + 1)$ for $(*) \frac{d\mathcal{I}_{\beta}}{d\beta} = 0$ $\stackrel{(*)}{=} \ln 2 + (1 - \delta) \ln(1 - \delta) + \delta \ln \delta$

 $\left(\delta = \frac{1}{n} |\{i : \xi_i^{(1)} \neq \xi_i^{(2)}\}|\right)$

Channel capacity of BSC

 $\xi^{(2)}$





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ASC selects optimal (true) number of clusters Experimental Setting:

5 Gaussians, n=10000, d=2, k^{max}=10



Role-Based Access Control

- Given: Binary user permission matrix ^{change group web-page}
- Discretional supervise master thesis
 Access-Control: use coffee machine



NAMES OF STREET

Direct Assignments of users to permissions

 Role-Based Access Control (RBAC): Permissions are granted via roles



Role-Mining for RBAC

- Role-Mining: Given a user-permission assignment matrix X, find a set of roles U and assignments Z such that
 X \approx U \otimes Z
- Multi Assignment Clustering: generative approach including noise model, inference with DA



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Synthetic Data: Parameter Accuracy vs. Approximation Capacity



ASC ranking of model variants complies with ranking according to ground truth.

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Real-World Data: Prediction Error complies with Approximation Capacity

• Generalization: Can roles predict permissions of new

users?

- 1. Use few permissions (20%)to determine role setImage: Second second
- to determine role set 2. Predict hidden/missing permissions (80%). Centroids with maximal
- Centroids with maximal capacity yield minimal generalization error



Denoising Binary Matrices by truncated SVD

 $X_{5} = U_{5}S_{5}V_{5}$ Boolean matrix with 40% random entries Biede=1; White=0 6N X = USVRounding as approximation $g(X_k) = round(X_k)$

20 26

46 60

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Maximum of approximation capacity selects optimal rank k

Integrate over variations of the signal matrix U.



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Conclusion

- Quantization: Noise quantizes hypothesis classes => symbols
- These symbols can be used for coding!
- Optimal error free coding scheme determines
 approximation capacity of a cost function.
- \Rightarrow Bounds for robust optimization.
- ⇒Quantization of hypothesis class measures structure specific information in data.

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X

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Future Work

- Generalization: replace approximation sets based on cost functions by smoothed outputs of algorithms ("smoothed generalization")
- Model reduction in dynamical systems: quantize sets of ODEs or PDEs (systems biology)
- Relate statistical complexity, i.e. the approximation capacity, to algorithmic or computational complexity.